

The GNU Name System

Christian Grothoff

Technische Universität München

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“Never doubt your ability to change the world.” –Glenn Greenwald

The Internet

Virtually all Internet protocols are broken:

Ethernet MAC spoofing, cleartext

IP IP spoofing, cleartext

BGP AS hijacking, cleartext

DNS cache poisoning, cleartext

DNSSEC cleartext, often no end-to-end authentication

TLS 100 CAs can certify anybody for anything

HTTP too chatty, complex, slow

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Rule 1 for the GUNet: Encrypt everything.

Encryption to the Rescue?

- ▶ Existing Internet PKIs are easily controlled:
 - ▶ DNSSEC root certificate
 - ▶ X.509 CAs (HTTPS certificates)
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The GNU Name System¹

Properties of GNS

- ▶ Decentralized name system with secure memorable names
- ▶ Delegation used to achieve transitivity
- ▶ Achieves query and response privacy
- ▶ Provides alternative public key infrastructure
- ▶ Interoperable with DNS

¹Joint work with Martin Schanzenbach and Matthias Wachs

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- ▶ A of **pixel.net.in.tum.de** is `131.159.20.32`

Zone Management: like in DNS


gnunet-setup


General Network Transports File Sharing Namestore **GNS**

Editing zone APISQDP7A126P06VV60535PDT50B9L12NK6QP64IE8KNC6E807G0 

Preferred zone name (PSEU):

Master Zone Private Zone Shorten Zone

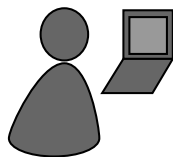


 Save As

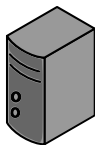
Name	Type	Value	Expiration	Public
<new name>				
+ >	<new record>			
	MX	5,mail.+	end of time	<input checked="" type="checkbox"/>
priv >	<new record>			
	PKEY	3IQ1TG601GUBVO55C0J087OEFB8N3DBJQ4L9SBI8PFLR8UKCVGHG	end of time	<input type="checkbox"/>
heise >	<new record>			
	LEHO	heise.de	end of time	<input checked="" type="checkbox"/>
	AAAA	2a02:2e0:3fe:100::8	end of time	<input checked="" type="checkbox"/>
	A	193.99.144.80	end of time	<input checked="" type="checkbox"/>
home >	<new record>			
大学 >	<new record>			
short >	<new record>			
mail >	<new record>			
homepage >	<new record>			
fcfs >	<new record>			
www >	<new record>			

[Welcome to gnunet-setup.](#)


Name resolution in GNS



Bob



Bob's webserver

Local Zone: K_{pub}^{Bob}		
www	A	5.6.7.8
		

- ▶ Bob can locally reach his webserver via **www.gnu**

Secure introduction



TUM

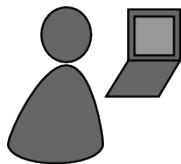


Bob Builder, Ph.D.

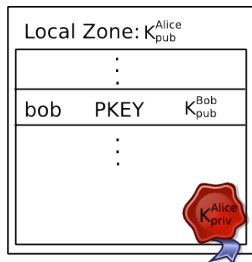
Address: Country, Street Name 23
Phone: 555-12345
Mobile: 666-54321
Mail: bob@H2R84L4JIL3G5C.zkey

- ▶ Bob gives his public key to his **friends**, possibly via QR code

Delegation

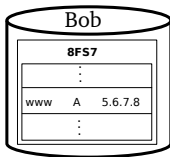
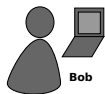


Alice

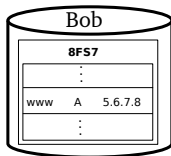
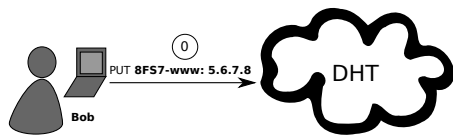


- ▶ Alice learns Bob's public key
- ▶ Alice creates delegation to zone K_{pub}^{Bob} under label **bob**
- ▶ Alice can reach Bob's webserver via **www.bob.gnu**

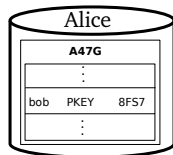
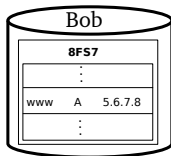
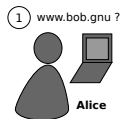
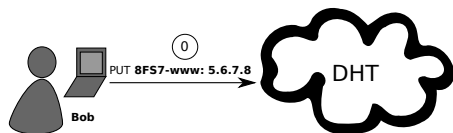
Name Resolution



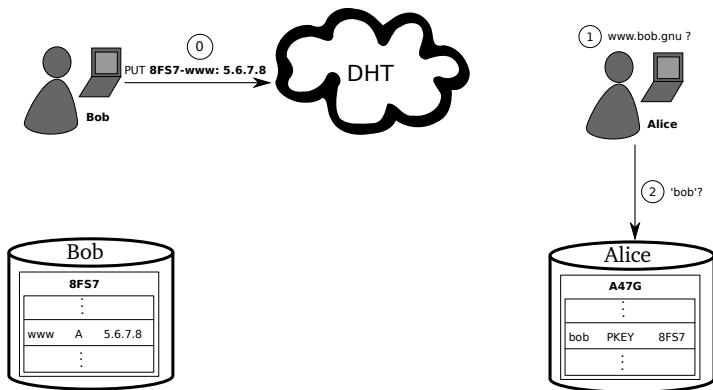
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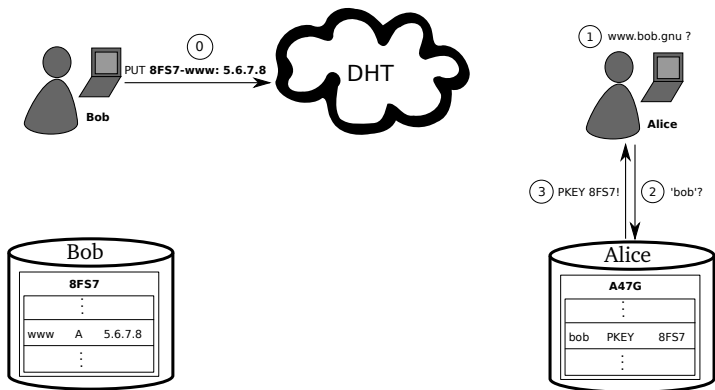
Name Resolution



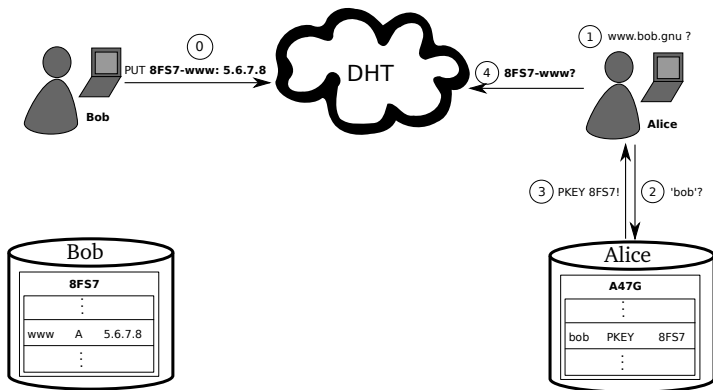
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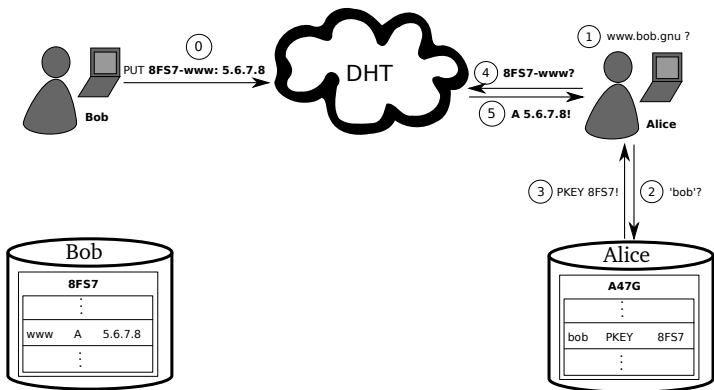
Name Resolution



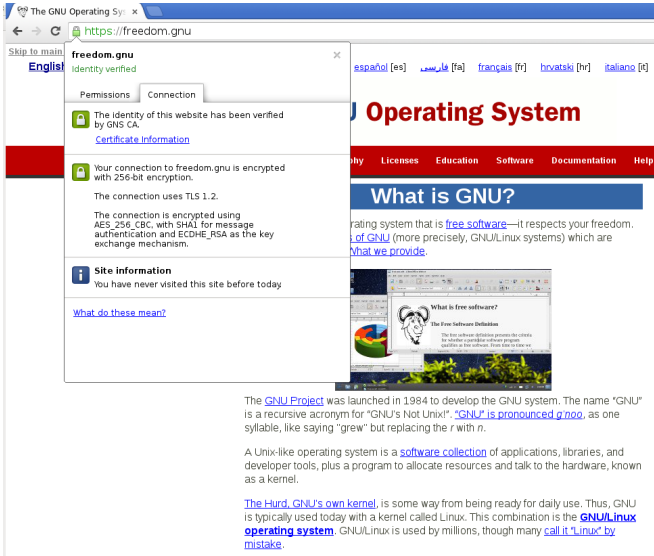
Name Resolution



Name Resolution



GNS as PKI (via DANE/TLSA)



The screenshot shows a web browser window with the address bar displaying `https://freedom.gnu`. A security warning dialog box is open, titled "freedom.gnu" with the subtext "identity verified". The dialog has two tabs: "Permissions" and "Connection".

Permissions

- The identity of this website has been verified by GNS CA. [Certificate Information](#)

Connection

- Your connection to freedom.gnu is encrypted with 256-bit encryption. The connection uses TLS 1.2. The connection is encrypted using AES_256_CBC, with SHA1 for message authentication and ECDHE_RSA as the key exchange mechanism.

Site information

- You have never visited this site before today. [What do these mean?](#)

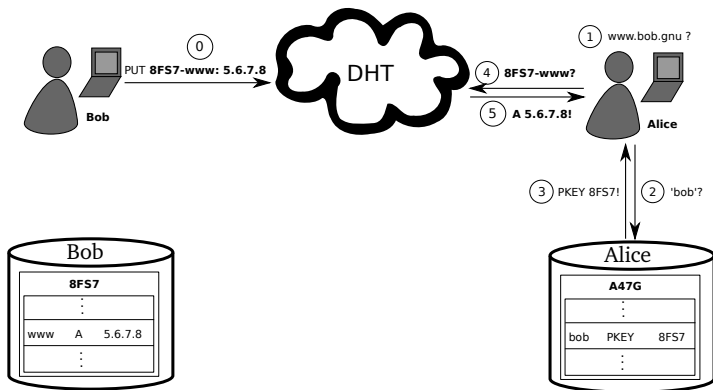
The background shows the GNU Operating System website. The main heading is "Operating System" in red. Below it is a navigation menu with links: "Why", "Licenses", "Education", "Software", "Documentation", and "Help". A blue banner reads "What is GNU?". Below the banner, the text says "operating system that is [free software](#)—it respects your freedom. of GNU (more precisely, GNU/Linux systems) which are [What we provide](#)." There is also a small image of a computer screen showing a "What is free software?" slide with a cartoon character.

The [GNU Project](#) was launched in 1984 to develop the GNU system. The name "GNU" is a recursive acronym for "GNU's Not Unix!". "GNU" is pronounced *g'noo*, as one syllable, like saying "grew" but replacing the *r* with *n*.

A Unix-like operating system is a [software collection](#) of applications, libraries, and developer tools, plus a program to allocate resources and talk to the hardware, known as a kernel.

[The Hurd, GNU's own kernel](#), is some way from being ready for daily use. Thus, GNU is typically used today with a kernel called Linux. This combination is the [GNU/Linux operating system](#). GNU/Linux is used by millions, though many [call it "linux" by mistake](#).

Privacy Issue: DHT



Query Privacy: Terminology

G generator in ECC curve, a point

n size of ECC group, $n := |G|$, n prime

x private ECC key of zone ($x \in \mathbb{Z}_n$)

P public key of zone, a point $P := xG$

l label for record in a zone ($l \in \mathbb{Z}_n$)

$R_{P,l}$ set of records for label l in zone P

$q_{P,l}$ query hash (hash code for DHT lookup)

$B_{P,l}$ block with encrypted information for label l
in zone P published in the DHT under $q_{P,l}$

Query Privacy: Cryptography

Publishing records $R_{P,I}$ as $B_{P,I}$ under key $q_{P,I}$

$$h := H(I, P) \quad (1)$$

$$d := h \cdot x \pmod n \quad (2)$$

$$B_{P,I} := S_d(E_{HKDF(I,P)}(R_{P,I})), dG \quad (3)$$

$$q_{P,I} := H(dG) \quad (4)$$

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Searching for records under label I in zone P

$$h := H(I, P) \quad (5)$$

$$q_{P,I} := H(hP) = H(hxG) = H(dG) \Rightarrow \text{obtain } B_{P,I} \quad (6)$$

$$R_{P,I} = D_{HKDF(I,P)}(B_{P,I}) \quad (7)$$

Revocation

Revocation Basics

- ▶ Revocation certificate (RC): message signed with private key
 - ▶ Peer receives new valid RC, floods to all neighbours
 - ▶ All peers store all valid RCs forever
- ⇒ Expensive operation ⇒ proof-of-work

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Revocation Magic

- ▶ Peers maybe offline during initial flood
 - ▶ Network might be temporarily partitioned
- ⇒ Need to reconcile revocation sets on connect

Whenever two peers establish a P2P connection, they must compute the set union of their RC sets!

The “.zkey” pTLD

- ▶ “LABELS.PKEY.zkey” format
 - ▶ PKEY is the public key of the zone
 - ▶ Works a bit like “.onion”
- ⇒ Globally unique identifiers!



NICKnames

- ▶ “alice.bob.carol.dave.gnu” is a bit long for Edward (“.gnu”)
- ▶ Also, we need to trust Bob, Carol and Dave (for each lookup)
- ▶ Finally, Alice would have liked to be called Krista (just Bob calls her Alice)

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- ▶ Eve learns the “NICK”, and GNS creates “krista.short.gnu”
- ▶ Memorable, short trust path in the future! TOFU!
- ▶ Krista better pick a reasonably unique NICK.

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- ▶ Records change
- ▶ Expiration time controls validity, like in DNS
- ▶ DHT propagation has higher delays, compared to DNS
- ▶ SHADOW is a flag in a record
- ▶ Shadow records are only valid if no other, non-expired record of the same type exists

Practical Concerns

- ▶ Name registration
- ▶ Support for browsing
- ▶ New record types
- ▶ Integration with applications
- ▶ State of the implementation

Registering a name in GNS

- ▶ Bob gives his PKEY to his **friends** via QR code
- ▶ or registers it at the **GNUnet fcfs** authority *pin.gnu* as "bob"
- ▶ → Bob's friends can resolve his records via **.petname.gnu*
- ▶ → or **.bob.pin.gnu*

From DNS to GNS

Names are not globally unique, but ...

... we need support for Virtual Hosting!

... we need support for SSL!

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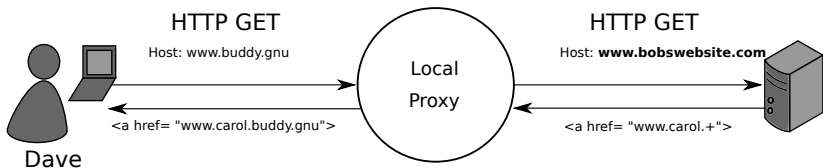
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Solution: Client Side SOCKS Proxy

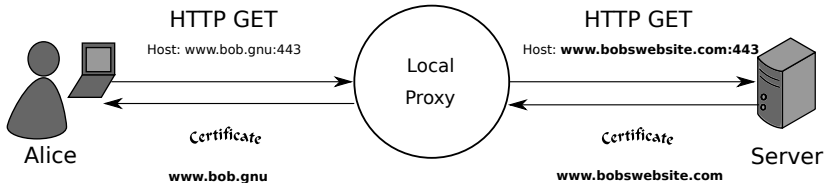
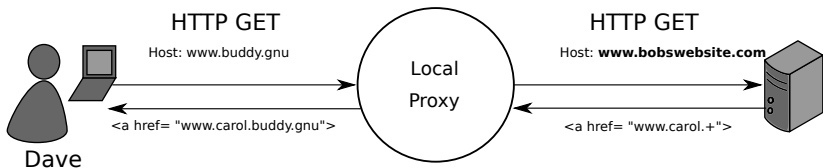
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Long-Term Vision

- ▶ Integration with browser and HTTP server
- ▶ HTTP server receives “GNS-Zone: PKEY” instead of “Hostname”
- ▶ HTTP client uses “TLSA” record of GNS, instead of “LEHO”

Relative Names

- ▶ GNS records can contain “.+”
- ▶ CNAME: “server1.+”
- ▶ MX: “mail.+”
- ▶ “.+” stands for “relative to current zone”

Supporting this for links in browsers would be nice, too.

New Record Types

- ▶ PKEY: delegate to another GNS zone
- ▶ NICK: preferred names for shortening
- ▶ LEHO: legacy hostname

New Record Types

- ▶ PKEY: delegate to another GNS zone
- ▶ NICK: preferred names for shortening
- ▶ LEHO: legacy hostname
- ▶ GNS2DNS: delegate to DNS
- ▶ VPN: peers hosting TCP/IP services
- ▶ PHONE: call users using `gnunet-conversation`
- ▶ BOX: proper support for TLSA (and SRV)

DNS Delegation

- ▶ Delegate to DNS using GNS2DNS records
- ▶ GNS2DNS record specifies:
 - ▶ Name of DNS resolver (i.e. “ns1.example.com” or “piratedns.”)
 - ▶ DNS domain to continue resolution in (i.e. “example.com” or “piratebay.org”)
- ▶ GNS will first resolve DNS resolver name to A/AAAA record
- ▶ GNS will then resolve “*left.of.gns2dns.example.com*” using DNS

VPN Delegation

- ▶ Delegates to GUNet VPN
- ▶ VPN record specifies:
 - ▶ Identity of hosting peer (no anonymity!)
 - ▶ Service identifier (hash code)
- ▶ GNS can map VPN record to A/AAAA record of `gnunet-vpn` tunnel

PHONE service

- ▶ PHONE record specifies:
 - ▶ Identity of hosting peer (no anonymity yet!)
 - ▶ Line number (to support multiple phones per peer)
- ▶ `gnunet-conversation` uses *reverse lookup* for caller ID

BOX records

- ▶ TLSA records in DNS are under a special name
- ▶ Performing a second lookup is **bad**

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- ▶ TLSA records in DNS are under a special name
- ▶ Performing a second lookup is **bad**
- ⇒ GNS BOX records include TLSA information under primary label!

Application Integration

- ▶ SOCKS proxy (`gnunet-gns-proxy`)
- ▶ NSS plugin
- ▶ DNS packet interception (`gnunet-dns-service`)
- ▶ GNS (C) API
- ▶ GNS (IPC) protocol
- ▶ GNS command-line tool

Current State

- ▶ GNS part of GUNet since 0.9.3
- ▶ Crypto changed to Curve25519 in 0.10.0
- ▶ Internationalized Domain Names are supported

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- ▶ GNS part of GUNet since 0.9.3
- ▶ Crypto changed to Curve25519 in 0.10.0
- ▶ Internationalized Domain Names are supported
- ▶ Installation is “non-trivial” (for your parents)
- ▶ SOCKS proxy is known to be problematic

Conclusion

- ▶ Decentralization is necessary
- ▶ Encryption requires a PKI
- ▶ GNS is a modern PKI designed for privacy
- ▶ Please consider adding GNS support to your code!

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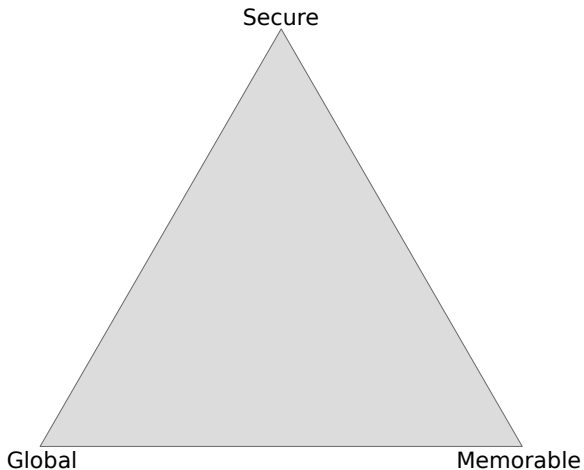


Do you have any questions?

References:

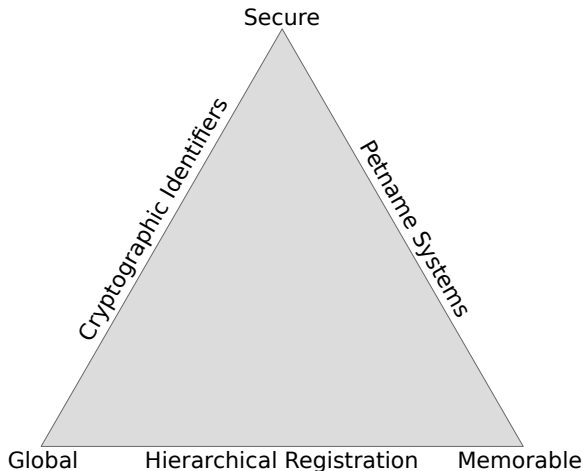
- ▶ Nathan Evans and Christian Grothoff. *R5N. Randomized Recursive Routing for Restricted-Route Networks*. **5th International Conference on Network and System Security**, 2011.
- ▶ Matthias Wachs, Martin Schanzenbach and Christian Grothoff. *On the Feasibility of a Censorship Resistant Decentralized Name System*. **6th International Symposium on Foundations & Practice of Security**, 2013.
- ▶ M. Schanzenbach *Design and Implementation of a Censorship Resistant and Fully Decentralized Name System*. **Master's Thesis (TUM)**, 2012.

Zooko's Triangle



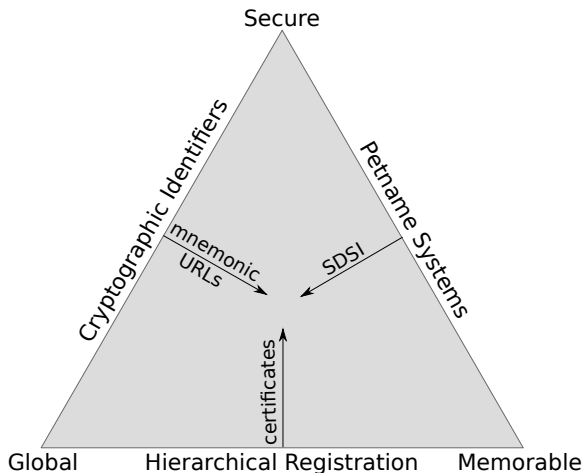
A name system can only fulfill **two!**

Zooko's Triangle



DNS, “.onion” IDs and `/etc/hosts/` are representative designs.

Zooko's Triangle



DNSSEC security is broken by design (adversary model!)

Namecoin

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- ▶ Memorable:

Namecoin

- ▶ Memorable: Check
- ▶ Global:

Namecoin

- ▶ Memorable: Check
- ▶ Global: Check
- ▶ Secure:

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- ▶ Secure: different adversary model!
- ⇒ Availability of names (registration rate) is restricted
- ⇒ Adversary must not have 51% compute power